

Wed., 05/16/12:

HW 6-7: some example numbers: time in HH:MM

$$N \times N = 4096 \times 4096:$$

4 cores		1 node
1 process		1:37
2		0:55
4		0:29

8 ranks		1 node
1		0:57
2		0:39
4		0:32
8		0:16

## Ch. 10 Design of Parallel Algorithms

Part 1: Jacobi method = iterative method for  $Ax = b$

→ idea of matrix-free method

→ We did already CG, which is much better mathematically

Issue is that Jacobi needs millions of iterations for moderately large  $N$ , such as 8,192.  $\Rightarrow$  There is no way to run this to convergence, even on many nodes

$\Rightarrow$  It is always more important to use a good numerical method than a parallel one!

Part 2 of Ch.10: Parallel sorting , really: introduction of some new MPI commands

Introduce MPI\_Alltoall and MPI\_Alltoallv ,  
then also know about MPI\_Gatherv , etc.

→ Be aware of Appendix A .

Example : here a little different from Pacheco .

Assume  $p=4$  processes, problem: Have a total of 100 integers  $1, 2, \dots, 100$  that are given in local vectors "unsorted" of length 25 across all processes .

Want to obtain:  $1, 2, \dots, 25$  on Proc 0 ,  $26, \dots, 50$  on Proc 1 ,

etc. , in local vector "sorted" .

Critical: Never assemble any vector of length  $> 25$  anywhere!

Simplified from Pacheco : We do not allow for duplicate (or missing) numbers (like  $1, 1, 3, 4, \dots$ )  
⇒ We can allocate all vectors once, since we know their lengths .

Ex.:

Proc 0 : unsorted = [  $\underbrace{2}_{\text{to } P0}, \underbrace{26}_{\text{to } P1}, \underbrace{51, 52, \dots, 72}_{\text{to } P2}, \underbrace{100}_{\text{to } P3}$  ]

Proc 1 : [  $\underbrace{1, 2}_{\text{to } P0}, \underbrace{73}_{\text{to } P1}, \underbrace{76, \dots, 97}_{\text{to } P3} ]$

.

Need communication, so that we arrive at

Proc 0 : sorted = [ 2, 1, ...  
| : [ 26, 27, ..., ]  
2 [ 51, 52, ...  
3 [ 100, 76, 77, ... ]

$\Rightarrow$  "sorted" will not be sorted after the communication, even if unsorted on each process are sorted locally.

MPJ-All to all {

above example  
 $\downarrow$

void \* send\_buf  
int send\_count  
MPI\_Datatype send\_type

unsorted  
①

void \* recv\_buf  
int recv\_count  
MPI\_Datatype recv\_type

MPI\_INT  
sorted  
②

MPI\_Datatype  
MPI\_Comm  
MPI\_Comm ) MPI\_COMM\_WORLD

① This needs to be same number to each process! That is just not the case  
But we can determine the needed numbers from the data locally.

② We cannot determine this?

For ① and ②, need variable version  
=

MPI\_Neighbor (

void	* send - buffer
int	* send counts
int	* send - displacements
MPI_Datatype	send type
void	* recv - buffer
int	* recv - counts
int	* recv - displacements
MPI_Datatype	recv - type
MPI_Comm	comm )

only output is recv - buffer; everything else is

How long are vectors send - counts, send - displacements, recv - counts, recv - displacements? np

=> need to allocate dynamically

Write code :

```
int *unsorted, *send_counts, *send_displacements  
int *sorted, *recv_counts, *recv_displacements  
/* np divides n = 100 assumed ! */  

$$l_n = \frac{n}{np}$$
  
unsorted = allocate_int_vector( l_n )  
sorted = ...  
send_counts = allocate_int_vector( np )  
recv_counts = ...  
send_displacements = ...  
recv_displacements = ...
```

Perform local sorting  $\Rightarrow$  now have unsorted  
sorted, as in example above.

```
for (i=0; i < np; i++) send_counts[i] = 0
```

```
for (j=0; j < l_n; j++)
```

```
(send_counts[(unsorted[j]-1)/l_n])++
```

```
send_displacements[0] = 0
```

```
for (i=1; i < np; i++)
```

```
send_displacements[i] =
```

```
send_displacements[i-1] + send_counts[i-1]
```

② We need to get the recv\_counts ?

```
MPI_Alltoall( send_counts, 1, MPI_INT,  
               recv_counts, 1, MPI_INT,  
               MPI_COMM_WORLD )
```

recv\_displacements [0]

for ( $i = 0$ ;  $i < np$ ;  $i++$ )

recv\_displacements [ $i$ ] =

recv\_displacements [ $i - 1$ ] + recv\_counts [ $i - 1$ ]

MPI\_Alltoallv (unrooted, sendcounts,  
send\_displacements, MPI\_INT,  
sorted, recv\_count, recv\_displacements,  
MPI\_INT, MPI\_COMM\_WORLD )

Perform local search on "sorted" !

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Bigger message: There are " $\wedge$ " version of  
other collective commands like MPI\_Gatherv

MPI\_Scatterv, etc.

See man page for syntax

See also Appendix A, for instance to locate  
command that is nearly the right one,  
then see if other ones exist in that section  
that are better.

## Ch. 6 Derived Datatypes = Grouping Data

This will explain, why send-type can be different from recv-type.

Let  $A \in \mathbb{R}^{n \times n}$  be in 1-D column-oriented array storage (like A in power method)  
Let A be given on Process 0 and want to get  $l \cdot A \in \mathbb{R}^{n \times l \cdot n}$  on each process.

You can do this with

```
MPI_Scatter( A, 1, block_col_type,  
             l_A, (n*l_n), MPI_DOUBLE,  
             >0, MPI_COMM_WORLD )
```

Here, block-row-type is defined at run time

$\text{MPI_Type_vector}($  this example  $\rightarrow$   
    int count  $l \cdot n$   
    int blocklength  $n$   
    int stride 1

$\text{MPI_Datatype el\_type}$   $\text{MPI_DOUBLE}$

$\text{MPI_Datatype *new_mpi\_t})$  & block\_col\_type



`MPI_Type_vector ( n , l_n , n ,`

`MPI_DOUBLE, & block_row_type)`

each block needs to be consecutive

Above `block_col_type` has all data consecutive

$\Rightarrow$  could have done as 1 block of length  $n * l_n$

Or in fact of course, we could avoid derived datatype altogether, if data is consecutive!

But for MPI\_Scatter ( $A$ , 1, `block_row-type`)  
 the data is not consecutive and this derived datatype useful  $\rightarrow$  Unfortunately need to  
 fix the `block_row-type` as in Sec. 8.4.5  
 with MPI\_UB ?

Generally, regarding use of different send\_type  
 and recv\_type : the criterion is that their  
 type signatures have to match