

Wednesday, 12/21/11:

Representation of Matrices in C:

Perspective of Fortran = standard C perspective

double A[10][10] or

double A[n][n]

How are the data actually organized?

dynamic way:

double \*\*A

A = (double\*\*) malloc (n \* sizeof (double\*))

for (i = 0; i < n; i++)

A[i] = (double\*) malloc (n \* sizeof (double)) ← row

⇒ element A[i][j] = A<sub>ij</sub>.

⇒ the malloc's are all independent, so data are not consecutive in memory!

The same is true for A[n][n].

Also notice: C organizes matrices by row = row-major ordering. This is incompatible with Fortran and Matlab and thus also all major numerical libraries. More subtly, the

double pointer in A[n][n] is a problem, since no other language perceives arrays in this way. A[i][j] is simply not the same as

A[i, j].

Therefore, we 1-D arrays for all arrays including 2-D matrices or higher-dimensional arrays and then index yourself manually.

double \*A

A = (double\*) malloc(n\*n\*sizeof(double))

Use it as  $A[i+n*j] = A_{ij}$  (or  $A[i+n*j+n*n*k] = A_{ijk}$ )

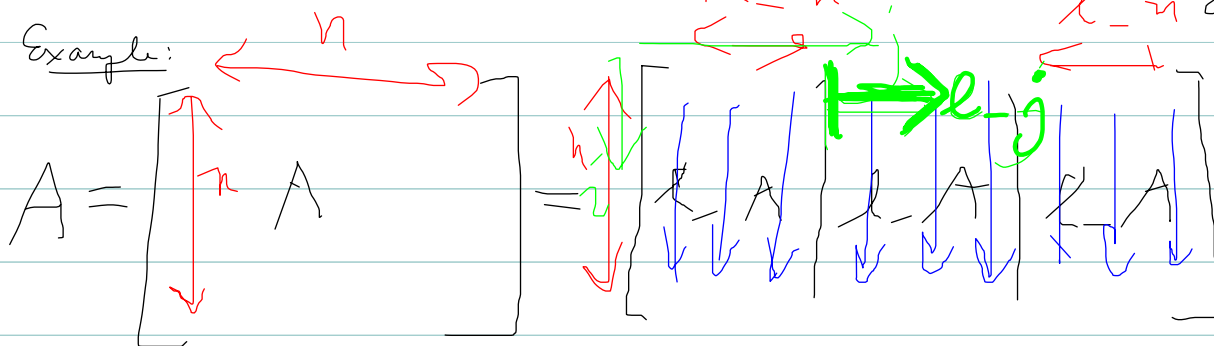
This is column-major ordering, since  $A[i+1+n*j]$  is next to  $A[i+n*j]$  in memory, which is  $A_{i+1,j}$  being next to  $A_{i,j}$ .

⇒ This way we are in complete control of where our data are!

This is vital to using MPI communication commands and to efficient

code.

Example:



We divide into blocks of columns to accommodate the column-major ordering,

since then also  $l-A$  are consecutive in memory.

⇒ Now we can do a single MPI\_Gather to collect all  $l-A$  into A on Process 0 for fast output (→ HW (a))

MPI\_Gather( $l-A$ ,  $n * l_n$ , MPI\_DOUBLE, A,  $n * l_n$ , MPI\_DOUBLE, 0, MPI\_COMM\_WORLD)

Clarify:

double \*l\_A

l\_A = allocate\_double\_vector( $n * l_n$ )

for ( $l_j < 0$ ;  $l_j < l_n$ ;  $l_j++$ )

for ( $i = 0$ ;  $i < n$ ;  $i++$ )

l\_A[ $i + n * l_j$ ] = A<sub>ij</sub>

$$j = l_j + id * l_n$$