IS 450/IS 650–
Data Communications and Networks

Course Review
Midterm Exam

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Midterm Exam

- When: Tuesday (10/27) 4:30pm - 6:30pm
- Where: In Class

- Closed book, Closed notes
- Computer Networks and the Internet (Chapter 1); Application Layer (Chapter 2) and Transport Layer (up to Chapter 3.3)

- Material for preparation:
  - Lecture Slides
  - Quizzes
  - Textbooks
    - Computer Networking: A Top Down Approach,
Course Overview

- **Computer Networks and the Internet (Chapter 1)**
  - Packet and circuit switching
  - End to End Delay

- **Application Layer (Chapter 2)**
  - Web, HTTP, FTP, SMTP, DNS, Peer-peer and Socket

- **Transport Layer (Chapter 3)**
  - Multiplexing & Demultiplexing
Chapter 1: Computer Networks & Internet

1.1 What is the Internet?
1.2 Network edge
1.3 Network access and physical media
1.4 Network core
1.5 Internet structure and ISPs
1.6 Delay and loss in packet-switched networks
1.7 Protocol layers, service models
1.8 History
Chapter 1: Roadmap

1.1 What is the Internet?
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1.8 History
The Network Core

- mesh of interconnected routers
- the fundamental question: how is data transferred through net?
  - circuit switching: dedicated circuit per call: telephone net
  - packet-switching: data sent thru net in discrete “chunks”
  - Forwarding table and routing protocols
Network Core: Circuit Switching

End-end resources reserved for “call”

- link bandwidth, switch capacity
- dedicated resources: no sharing
- circuit-like (guaranteed) performance
- call setup required
Network Core: Circuit Switching

Network resources (e.g., bandwidth) divided into “pieces”

- pieces allocated to calls
- resource piece *idle* if not used by owning call (*no sharing*)

- dividing link bandwidth into “pieces”
  - frequency division
  - time division
Circuit Switching: FDM and TDM

Example:
4 users

FDM
- Frequency
- Time

TDM
- Frequency
- Time
FDM vs TDM

- What are the tradeoffs?
  - Advantage and disadvantage of dividing frequency?
  - Advantage and disadvantage of dividing time?
Numerical example

How long does it take to send a file of 640,000 bits from host A to host B over a circuit-switched network?

- All links are 1.536 Mbps
- Each link uses TDM with 24 slots/sec
- 500 msec to establish end-to-end circuit

Let’s work it out!
Network Core: Packet Switching

each end-end data stream divided into packets
- user A, B packets share network resources
- each packet uses full link bandwidth
- resources used as needed

resource contention:
- aggregate resource demand can exceed amount available
  - Packets queue up
- store and forward: packets move one hop at a time
  - Node receives complete packet before forwarding

Bandwidth division into “pieces”
- Dedicated allocation
- Resource reservation
Packet Switching: Statistical Multiplexing

Sequence of A & B packets does not have fixed pattern, shared on demand **statistical multiplexing**.

TDM: each host gets same slot in revolving TDM frame.
Packet-switching: store-and-forward

- Takes $L/R$ seconds to transmit (push out) packet of $L$ bits on to link of $R$ bps
- Entire packet must arrive at router before it can be transmitted on next link: *store and forward*
- delay = $3L/R$ (assuming zero propagation delay)

**Example:**
- $L = 7.5$ Mbits
- $R = 1.5$ Mbps
- delay = 15 sec

more on delay shortly …
Packet-switched networks: forwarding

- **Goal:** move packets through routers from source to dest.
  - we’ll study several path selection (routing) algorithms (chap 4)

- datagram network:
  - *destination address* in packet determines next hop
  - routes may change during session
  - analogy: driving, asking directions

- virtual circuit network:
  - packet carries tag (virtual circuit ID), tag determines next hop
  - fixed path determined at *call setup time*, remains fixed thru call
  - *routers maintain per-call state*
  - (analogy: air trains in airports)
Compare

Thoughts on tradeoffs between packet switching and circuit switching?

Which one would you take?

Under what circumstances?

Why?
Packet switching versus Circuit switching

Packet switching allows more users to use network!

- problem: 1 Mbps link
- each user:
  - 100 kbps when “active”
  - active 10% of time
- circuit-switching:
  - 10 users
- packet switching (ps):
  - with 35 users, probability > 10 active users is less than 0.0004

Q: how did we get value 0.0004?
Get performance of circuit switching with 3 times more users in case of PS
Packet switching versus Circuit switching

Is packet switching a “slam dunk winner?”

- Great for absorbing bursty data from individual sources
  - resource sharing (due to diversity)
  - simpler, no call setup

- Excessive congestion: packet delay and loss
  - protocols needed for reliability, congestion control

- Q: How to provide circuit-like behavior?
  - bandwidth guarantees needed for audio/video apps
  - still unsolved (chapter 7)
Problem on Circuit and Packet switching

- Suppose users share a 15 Mbps link. Also suppose each user requires 1 Mbps when transmitting, but each user transmit only 10% time.

  a) When circuit switching is used, how many users can be supported?

  b) Suppose there are 30 users. Find the probability that any given time, exactly 20 users are transmitting simultaneously. (Hint: Use the binomial distribution)

- Solve this problem from Quiz 1
Binomial Probability Formula

\[ P(r) = ^nC_r p^r (1 - p)^{n-r} = \frac{n!}{r!(n-r)!} p^r q^{n-r} \]

for \( r = 0, 1, 2, \ldots, n \)

where

\( n \) = number of trials

\( r \) = number of successes among \( n \) trials

\( p \) = probability of success in any one trial

\( q \) = probability of failure in any one trial (\( q = 1 - p \))
Chapter 1: Roadmap

1.1 What is the Internet?
1.2 Network edge
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1.5 Internet structure and ISPs
1.6 Delay & loss in packet-switched networks
1.7 Protocol layers, service models
1.8 History
How do loss and delay occur?

packets *queue* in router buffers

- packet arrival rate to link exceeds output link capacity
- packets queue, wait for turn

packets queueing (delay)

packet being transmitted (delay)

free (available) buffers: arriving packets dropped (loss) if no free buffers
Four Sources of Packet Delay

1. nodal processing:
   - check bit errors
   - determine output link

2. queueing:
   - time waiting at output link for transmission
   - depends on congestion level of router
Delay in packet-switched networks

3. Transmission delay:
- \( L \) = packet length (bits)
- \( R \) = link bandwidth (bps)
- time to send bits into link = \( \frac{L}{R} \)

4. Propagation delay:
- \( d \) = length of physical link
- \( s \) = propagation speed in medium (~\( 2 \times 10^8 \) m/sec)
- propagation delay = \( \frac{d}{s} \)

Note: \( R \) and \( s \) are very different quantities!
Comparing Transmission & Propagation Delays

- Transmission delay
  - Amount of time required to push out a packet
  - Function of the packet’s length & transmission rate of the link
  - Nothing to do with the distance between the two routers

- Propagation delay
  - Time it takes a bit to propagate from one router to the next
  - Function of the distance between two routers and propagation speed
  - Nothing to do with the packets’ length or transmission rate
Nodal delay

\[ d_{\text{nodal}} = d_{\text{proc}} + d_{\text{queue}} + d_{\text{trans}} + d_{\text{prop}} \]

- \( d_{\text{proc}} \) = processing delay
  - typically a few microsecs or less
- \( d_{\text{queue}} \) = queuing delay
  - depends on congestion
- \( d_{\text{trans}} \) = transmission delay
  - \( = L/R \), significant for low-speed links
- \( d_{\text{prop}} \) = propagation delay
  - a few microsecs to hundreds of msecs
Chapter 2: Application layer

- 2.1 Principles of network applications
- 2.2 Web and HTTP
- 2.3 FTP
- 2.4 Electronic Mail
  - SMTP, POP3, IMAP
- 2.5 DNS
- 2.6 P2P file sharing
- 2.7 Socket programming with TCP
- 2.8 Socket programming with UDP
- 2.9 Building a Web server
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Application architectures

- Client-server
- Peer-to-peer (P2P)
- Hybrid of client-server and P2P
Client-server architecture

server:
- always-on host
- permanent IP address
- server farms for scaling
- data centers

clients:
- communicate with server
- may be intermittently connected
- may have dynamic IP addresses
- do not communicate directly with each other
Pure P2P architecture

- no always-on server
- arbitrary end systems directly communicate
- peers are intermittently connected and change IP addresses
- Major Challenges: ISP friendly, Security, Incentives
- example: Gnutella (peer-to-peer file sharing network)

Highly scalable but difficult to manage
Hybrid of client-server and P2P

**Skype**
- Internet telephony app
- Finding address of remote party: centralized server(s)
- Client-client connection is direct (not through server)

**Instant messaging**
- Chatting between two users is P2P
- Presence detection/location centralized:
  - User registers its IP address with central server when it comes online
  - User contacts central server to find IP addresses of buddies
Internet transport protocols services

TCP service:
- connection-oriented: setup required between client and server processes
- reliable transport between sending and receiving process
- flow control: sender won’t overwhelm receiver
- congestion control: throttle sender when network overloaded
- does not provide: timing, minimum bandwidth guarantees

UDP service:
- unreliable data transfer between sending and receiving process
- does not provide: connection setup, reliability, flow control, congestion control, timing, or bandwidth guarantee

Q: Why bother? Why is there a UDP?
### Internet apps: application, transport protocols

<table>
<thead>
<tr>
<th>Application</th>
<th>Application layer protocol</th>
<th>Underlying transport protocol</th>
</tr>
</thead>
<tbody>
<tr>
<td>e-mail</td>
<td>SMTP [RFC 2821]</td>
<td>TCP</td>
</tr>
<tr>
<td>remote terminal access</td>
<td>Telnet [RFC 854]</td>
<td>TCP</td>
</tr>
<tr>
<td>Web</td>
<td>HTTP [RFC 2616]</td>
<td>TCP</td>
</tr>
<tr>
<td>file transfer</td>
<td>FTP [RFC 959]</td>
<td>TCP</td>
</tr>
<tr>
<td>streaming multimedia</td>
<td>proprietary (e.g. RealNetworks)</td>
<td>TCP or UDP</td>
</tr>
<tr>
<td>Internet telephony</td>
<td>proprietary (e.g., Vonage,Dialpad)</td>
<td>typically UDP</td>
</tr>
</tbody>
</table>
Chapter 2: Application layer

- 2.1 Principles of network applications
  - app architectures
  - app requirements
- 2.2 Web and HTTP
- 2.4 Electronic Mail
  - SMTP, POP3, IMAP
- 2.5 DNS
- 2.6 P2P file sharing
- 2.7 Socket programming with TCP
- 2.8 Socket programming with UDP
- 2.9 Building a Web server
Web and HTTP

First some jargon

- **Web page** consists of objects
- Object can be HTML file, JPEG image, Java applet, audio file,...
- Web page consists of **base HTML-file** which includes several **referenced objects**
- Each object is addressable by a **URL**
- **Example URL:**

  
  www.someschool.edu/someDept/pic.gif

  
  **host name**  
  **path name**
HTTP overview

HTTP: hypertext transfer protocol
- Web’s application layer protocol
- client/server model
  - *client*: browser that requests, receives, “displays” Web objects
  - *server*: Web server sends objects in response to requests
- HTTP 1.0: RFC 1945
- HTTP 1.1: RFC 2068
HTTP overview (continued)

Uses TCP:
- client initiates TCP connection (creates socket) to server, port 80
- server accepts TCP connection from client
- HTTP messages (application-layer protocol messages) exchanged between browser (HTTP client) and Web server (HTTP server)
- TCP connection closed

HTTP is “stateless”
- server maintains no information about past client requests

Protocols that maintain “state” are complex!
- past history (state) must be maintained
- if server/client crashes, their views of “state” may be inconsistent, must be reconciled
HTTP connections

Nonpersistent HTTP
- At most one object is sent over a TCP connection
- HTTP/1.0 uses nonpersistent HTTP

Persistent HTTP
- Multiple objects can be sent over single TCP connection between client and server
- HTTP/1.1 uses persistent connections in default mode
Non-Persistent HTTP: Response time

Round Trip Time (RTT) = time to send a small packet to travel from client to server and back.

**Response time:**
- one RTT to initiate TCP connection
- one RTT for HTTP request and first few bytes of HTTP response to return
- file transmission time

total = 2RTT+ <file transmit time>
Persistent HTTP

Nonpersistent HTTP issues:
- requires 2 RTTs per object
- OS overhead for each TCP connection
- browsers often open parallel TCP connections to fetch referenced objects

Persistent HTTP
- server leaves connection open after sending response
- subsequent HTTP messages between same client/server sent over open connection

Persistent without pipelining:
- client issues new request only when previous response has been received
- one RTT for each referenced object

Persistent with pipelining:
- default in HTTP/1.1
- client sends requests as soon as it encounters a referenced object
- as little as one RTT for all the referenced objects
DNS: Domain Name System

People: many identifiers:
  - SSN, name, passport #

Internet hosts, routers:
  - IP address (32 bit) - used for addressing datagrams
  - “name”, e.g., www.yahoo.com - used by humans

Q: map between IP addresses and name?

Domain Name System:

- *distributed database* implemented in hierarchy of many *name servers*

- *application-layer protocol* host, routers, name servers to communicate to *resolve* names (address/name translation)
  - note: core Internet function, implemented as application-layer protocol
  - complexity at network’s “edge”
DNS

DNS services

- Hostname to IP address translation
- Host aliasing
  - Canonical and alias names
  - Alias: enterprise.com or www.enterprise.com
    Canonical: relay1.west coast.enterprise.com
- Load distribution
  - Replicated Web servers: set of IP addresses for one canonical name

Why not centralize DNS?

- single point of failure
- traffic volume
- distant centralized database
- Maintenance doesn’t scale!
Distributed, Hierarchical Database

Client wants IP for www.amazon.com; 1st approx:
- Client queries a root server to find .com DNS server
- Client queries com DNS server to get amazon.com DNS server
- Client queries amazon.com DNS server to get IP address for www.amazon.com
DNS records

**DNS:** distributed db storing resource records (RR)

RR format: \((name, value, type, ttl)\)

- **Type=A**
  - name is hostname
  - value is IP address

- **Type=NS**
  - name is domain (e.g. foo.com)
  - value is hostname of authoritative name server for this domain (e.g. dns.foo.com)

- **Type=CNAME**
  - name is alias name for some “canonical” (the real) name
    - www.ibm.com is really servereast.backup2.ibm.com
  - value is canonical name

- **Type=MX**
  - value is name of mailserver associated with name (e.g. foo.com, mail.bar.foo.com, MX)
Example: just created startup “Network Utopia”
Register name networkutopia.com at a registrar (e.g., Network Solutions)
  - Need to provide registrar with names and IP addresses of your authoritative name server (primary and secondary)
  - Registrar inserts two RRs into the com TLD server:
    - (networkutopia.com, dns1.networkutopia.com, NS)
    - (dns1.networkutopia.com, 212.212.212.1, A)

Put in authoritative server Type A record for www.networkutopia.com and Type MX record for mail.networkutopia.com

How do people get the IP address of your Web site?
Chapter 3 Outline

- 3.1 Transport-layer services
- 3.2 Multiplexing and demultiplexing
- 3.3 Connectionless transport: UDP
- 3.4 Principles of reliable data transfer

- 3.5 Connection-oriented transport: TCP
  - segment structure
  - reliable data transfer
  - flow control
  - connection management
- 3.6 Principles of congestion control
- 3.7 TCP congestion control
Transport Layer (Chapter 3)

- Transport Layer
- Multiplexing / Demultiplexing
## Multiplexing/demultiplexing

### Demultiplexing at rcv host:
Delivering received segments to correct socket

### Multiplexing at send host:
Gathering data from multiple sockets, enveloping data with header (later used for demultiplexing)

<table>
<thead>
<tr>
<th>process</th>
<th>socket</th>
</tr>
</thead>
<tbody>
<tr>
<td>host 1</td>
<td>P3</td>
</tr>
<tr>
<td>host 2</td>
<td>P1</td>
</tr>
<tr>
<td>host 3</td>
<td>P4</td>
</tr>
</tbody>
</table>

One HTTP process, one FTP process, one Telnet process
More than one socket, each socket has unique identifier
How demultiplexing works

- host receives IP datagrams
  - each datagram has source IP address, destination IP address
  - each datagram carries 1 transport-layer segment
  - each segment has source, destination port number

- host uses IP addresses & port numbers to direct segment to appropriate socket

Analogous to car rentals at airports
Shuttles MUX passengers and take them To rental office -- DeMUX to diff cars
Connectionless demultiplexing

- Create sockets with port numbers:
  
  ```java
  DatagramSocket mySocket1 = new DatagramSocket(99111);
  DatagramSocket mySocket2 = new DatagramSocket(99222);
  ```

- UDP socket fully identified by two-tuple:
  
  `dest IP address, dest port number`

- When host receives UDP segment:
  
  - checks destination port number in segment
  - directs UDP segment to socket with that port number

- IP datagrams with different source IP addresses and/or source port numbers directed to same socket
DatagramSocket serverSocket = new DatagramSocket(6428);
Connection-oriented demux

- TCP socket identified by 4-tuple:
  - source IP address
  - source port number
  - dest IP address
  - dest port number

- recv host uses all four values to direct segment to appropriate socket

- Server host may support many simultaneous TCP sockets:
  - each socket identified by its own 4-tuple

- Web servers have different sockets for each connecting client:
  - non-persistent HTTP will have different socket for each request
Connection-oriented demux (cont)

Client IP: A
- SP: 9157
- DP: 80
- S-IP: A
- D-IP: C

Server IP: C
- SP: 5775
- DP: 80
- S-IP: B
- D-IP: C

Client IP: B
Connection-oriented demux:
Threaded Web Server

Client IP: A

Server IP: C

Client IP: B
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- 3.7 TCP congestion control
UDP: User Datagram Protocol [RFC 768]

- “no frills,” “bare bones” Internet transport protocol
- “best effort” service, UDP segments may be:
  - lost
  - delivered out of order to app
- connectionless:
  - no handshaking between UDP sender, receiver
  - each UDP segment handled independently of others

Why is there a UDP?
- no connection establishment (which can add delay)
- simple: no connection state at sender, receiver
- small segment header
- no congestion control: UDP can blast away as fast as desired
Tentative Midterm Exam Structure

- Short Multiple Choice Questions = 15 points
- Chapter 1: (packet/circuit switching, delay calculations etc.) = 30 points
- Chapter 2: (HTTP, Email, DNS etc.) = 30 points
- Chapter 3: (transport layer) = 15 points
- General Concepts: = 10 points

= 100 points

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Special office hours for IS 450/650 Midterm Exam:

Tuesday, Oct 27: 11am to 1:00pm at my office (ITE 421)

Good Luck!